

THE CONVERSE ENVELOPE THEOREM

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paper: [arXiv.org/abs/1909.11219](https://arxiv.org/abs/1909.11219)

Envelope theorem: optimal decision-making \implies \boxtimes formula.

Textbook intuition: \boxtimes formula \iff FOC.

Modern envelope theorem of MS02:* almost no assumptions.

\hookrightarrow FOC ill-defined, so need different intuition.

My theorem: with almost no assumptions,
 \boxtimes formula equivalent to generalised FOC.

– an envelope theorem: FOC \implies \boxtimes

– *a converse*: $\boxtimes \implies$ FOC.

Application to mechanism design.

*Milgrom, P., & Segal, I. (2002). Envelope theorems for arbitrary choice sets. *Econometrica*, 70(2), 583–601. doi:10.1111/1468-0262.00296

Setting

Agent chooses action x from a set \mathcal{X} .

Objective $f(x, t)$, where $t \in [0, 1]$ is a parameter.

No assumptions on \mathcal{X} , almost none on f :

(1) $f(x, \cdot)$ is differentiable for each $x \in \mathcal{X}$

(2) $f(x, \cdot)$ is ‘not too erratic’.

Decision rule: a map $X : [0, 1] \rightarrow \mathcal{X}$.

Associated value function: $V_X(t) := f(X(t), t)$.

Envelope theorem

X satisfies the \boxtimes formula iff

$$V_X(t) = V_X(0) + \int_0^t f_2(X(s), s) ds \quad \text{for every } t \in [0, 1].$$

Equivalently: V_X is absolutely continuous and

$$V'_X(t) = f_2(X(t), t) \quad \text{for a.e. } t \in (0, 1).$$

X is optimal iff for every t , $X(t)$ maximises $f(\cdot, t)$.

Modern envelope theorem (MS02).

Any optimal decision rule satisfies the \boxtimes formula.

Textbook intuition

Differentiation identity for $V_X(t) := f(X(t), t)$:

$$V'_X(t) = \underbrace{\frac{d}{dm} f(X(t+m), t) \Big|_{m=0}}_{\text{'indirect effect'}} + \underbrace{f_2(X(t), t)}_{\text{'direct effect'}}.$$

$$\begin{aligned} V'_X(t) &= \text{direct effect} && (\boxtimes \text{ formula}) \\ \iff \text{indirect effect} &= 0 && (\text{FOC}). \end{aligned}$$

Problem: 'indirect effect' (hence FOC) ill-defined!

- $f(\cdot, t)$ & X need not be differentiable.
- actions \mathcal{X} need have no convex or topological structure.

The outer first-order condition

Disjuncture: in general, \boxtimes formula $\not\leftrightarrow$ FOC.

- one solution: add strong ‘classical’ assumptions.
- my solution: find the correct FOC!

Decision rule X satisfies the outer FOC iff

$$\frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0} = 0 \quad \text{for all } r, t \in (0, 1).$$

‘Integrated’ version of classical FOC.

- always well-defined & necessary for optimality
- equiv’nt to classical FOC when latter well-defined.

Theorem

Envelope theorem & converse.

For a decision rule $X : [0, 1] \rightarrow \mathcal{X}$, the following are equivalent:

- (1) X satisfies the oFOC,
and $V_X(t) := f(X(t), t)$ is absolutely continuous.
- (2) X satisfies the \boxtimes formula

$$V_X(t) = V_X(0) + \int_0^t f_2(X(s), s) ds \quad \text{for every } t \in [0, 1].$$

Application: environment

Agent with preferences $f(y, p, t)$ over physical outcome $y \in \mathcal{Y}$ and payment $p \in \mathbf{R}$.

- \mathcal{Y} partially ordered
- type $t \in [0, 1]$ is agent's private info
- assume single-crossing.

A *physical allocation* is $Y : [0, 1] \rightarrow \mathcal{Y}$.

Y is *implementable* iff \exists payment rule $P : [0, 1] \rightarrow \mathbf{R}$
s.t. (Y, P) is incentive-compatible.

Application: goal

Classical result: implementable \iff increasing.

‘ \Leftarrow ’ is the substantial part. Versions:

	literature	this paper
outcomes \mathcal{Y}	$\subseteq \mathbf{R}$	general
preferences f	quasi-linear	general.

Application: result

Implementability theorem. Under regularity assumptions, any increasing physical allocation is implementable.

Argument:

- fix an increasing physical allocation $Y : [0, 1] \rightarrow \mathcal{Y}$
- choose a payment rule P so that \boxtimes formula holds
- then by *converse envelope theorem*, oFOC holds
 \iff mechanism (Y, P) is locally IC.
- finally, local IC \implies global IC by single-crossing.

Application: example

Monopolist selling information.

Physical allocations \mathcal{Y} :

distributions of posterior beliefs, ordered by Blackwell.

By the implementability theorem, any information allocation that gives higher types Blackwell-better signals can be implemented.

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Refresher: absolute continuity

Lipschitz \implies AC \implies uniformly continuous
& diff'able a.e.

Lebesgue's fundamental theorem of calculus.

For $\phi : [0, 1] \rightarrow \mathbf{R}$, the following are equivalent:

- (1) ϕ is AC.
- (2) ϕ is diff'able a.e. & $\phi(t) = \phi(0) + \int_0^t \phi' \quad \forall t \in [0, 1]$.

Refresher: uniform integrability

Family $\{\phi_i\}_{i \in I}$ of functions $[0, 1] \rightarrow \mathbf{R}$ is *uniformly integrable* iff for any $\varepsilon > 0$, $\exists \delta > 0$ s.t.

$$\sup_{i \in I} \int_T |\phi_i| < \varepsilon \quad \text{for any } T \subseteq [0, 1] \text{ of measure } < \delta.$$

Vitali convergence theorem.

If a sequence $\{\phi_n\}_{n \in \mathbf{N}}$ of functions $[0, 1] \rightarrow \mathbf{R}$ is UI & a.e. convergent, then $\lim_{n \rightarrow \infty} \phi_n$ is integrable and

$$\lim_{n \rightarrow \infty} \int_r^t \phi_n = \int_r^t \lim_{n \rightarrow \infty} \phi_n \quad \text{for any } r, t.$$

bounded \implies dominated by integrable f'n \implies UI.

Setting

Agent chooses action x from a set \mathcal{X} .

Objective $f(x, t)$, where $t \in [0, 1]$ is a parameter.

No assumptions on \mathcal{X} , almost none on f :

- (1) $f(x, \cdot)$ is differentiable for each $x \in \mathcal{X}$
- (2) $f(x, \cdot)$ is ‘not too erratic’.

Absolute equi-continuity

A family $\{\phi_x\}_{x \in \mathcal{X}}$ of functions $[0, 1] \rightarrow \mathbf{R}$ is *absolutely equi-continuous (AEC)* iff the family

$$\left\{ t \mapsto \sup_{x \in \mathcal{X}} \left| \frac{\phi_x(t+m) - \phi_x(t)}{m} \right| \right\}_{m>0}$$

is uniformly integrable.

(properties: slide 45)

‘ $f(x, \cdot)$ not too erratic’: $\{f(x, \cdot)\}_{x \in \mathcal{X}}$ is AEC.

- a sufficient condition (maintained by MS02):
 - $f(x, \cdot)$ absolutely continuous for each $x \in \mathcal{X}$, and
 - $t \mapsto \sup_{x \in \mathcal{X}} |f_2(x, t)|$ dominated by an integrable f’n.
- a stronger sufficient condition: f_2 bounded.

Setting

Agent chooses action x from a set \mathcal{X} .

Objective $f(x, t)$, where $t \in [0, 1]$ is a parameter.

No assumptions on \mathcal{X} , almost none on f :

(1) $f(x, \cdot)$ is differentiable for each $x \in \mathcal{X}$

(2) $f(x, \cdot)$ is ‘not too erratic’.

Decision rule: a map $X : [0, 1] \rightarrow \mathcal{X}$.

Associated value function: $V_X(t) := f(X(t), t)$.

Envelope theorem

X satisfies the \boxtimes formula iff

$$V_X(t) = V_X(0) + \int_0^t f_2(X(s), s) ds \quad \text{for every } t \in [0, 1].$$

X is optimal iff for every t , $X(t)$ maximises $f(\cdot, t)$.

Modern envelope theorem (MS02).[†]

Any optimal decision rule satisfies the \boxtimes formula.

[†]Really a slight refinement of MS02.

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Classical assumptions

Classical assumptions:

- \mathcal{X} is a convex subset of \mathbf{R}^n
- action derivative f_1 exists & is bounded
- only Lipschitz continuous decision rules X are considered.

(Bad for applications. Especially the Lipschitz restriction!)

Given a Lipschitz decision rule $X : [0, 1] \rightarrow \mathcal{X}$,

- type t can ‘mimic’ $t + m$ by choosing $X(t + m)$
- ‘mimicking payoff’ $m \mapsto f(X(t + m), t)$ is diff’able a.e.

Classical FOC: $\left. \frac{d}{dm} f(X(t + m), t) \right|_{m=0} = 0$ for a.e. t .

Classical envelope theorem and converse

Classical envelope theorem and converse.

Under the classical assumptions, classical FOC \iff \boxtimes formula.

Proof idea: Differentiation identity for $V_X(t) := f(X(t), t)$:

$$V'_X(t) = \underbrace{\frac{d}{dm} f(X(t+m), t) \Big|_{m=0}}_{\text{'indirect effect'}} + \underbrace{f_2(X(t), t)}_{\text{'direct effect'}}.$$

$$\begin{aligned} & V'_X(t) = \text{direct effect} \quad \text{a.e.} \quad (\boxtimes \text{ formula}) \\ \iff & \text{indirect effect} = 0 \quad \text{a.e.} \quad (\text{classical FOC}). \end{aligned}$$

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Trouble

Absent classical assumptions,

‘mimicking payoff’ $m \mapsto f(X(t+m), t)$ need not be diff’able.

- $f(\cdot, t)$ & X need not be diff’able
- actions \mathcal{X} need have no convex or topological structure.

\implies classical FOC ill-defined.

The outer first-order condition

Decision rule X satisfies the outer FOC iff

$$\frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0} = 0 \quad \text{for all } r, t \in (0, 1).$$

Motivation: given decision rule $X : [0, 1] \rightarrow \mathcal{X}$,

- type s can ‘mimic’ $s + m$ by choosing $X(s + m)$.
- oFOC: if types $s \in [r, t]$ do this,
it’s collectively unprofitable (to first order).

Housekeeping I

Housekeeping lemma. Under the classical assump'ns,
oFOC \iff classical FOC.

Sketch proof. Fix a Lipschitz $X : [0, 1] \rightarrow \mathcal{X}$.

Classical assump'ns & Vitali convergence theorem:

$$\underbrace{\frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0}}_{\substack{= 0 \text{ for all } r, t \\ \text{iff oFOC holds}}} = \underbrace{\int_r^t \frac{d}{dm} f(X(s+m), s) \Big|_{m=0} ds}_{\substack{= 0 \text{ for all } r, t \\ \text{iff classical FOC holds}}}.$$



Housekeeping II

Necessity lemma. Any optimal decision rule X satisfies the oFOC & has $V_X(t) := f(X(t), t)$ AC.

(sketch proof: slide 46)

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Envelope theorem & converse.

For a decision rule $X : [0, 1] \rightarrow \mathcal{X}$, the following are equivalent:

(1) X satisfies the oFOC

$$\frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0} = 0 \quad \text{for all } r, t \in (0, 1),$$

and $V_X(t) := f(X(t), t)$ is absolutely continuous.

(2) X satisfies the \boxtimes formula

$$V_X(t) = V_X(0) + \int_0^t f_2(X(s), s) ds \quad \text{for every } t \in [0, 1].$$

Main theorem

Envelope theorem & converse. For $X : [0, 1] \rightarrow \mathcal{X}$, TFAE:

- (1) X satisfies the oFOC, & $V_X(t) := f(X(t), t)$ is AC.
 - (2) X satisfies the \boxtimes formula.
-

\implies : an envelope theorem.

Implies the MS02 & classical envelope theorems.

\impliedby : *converse* envelope theorem.

Implies the classical converse envelope theorem.

The identity lemma

Proof in classical case relied on differentiation identity

$$V'_X(s) = \underbrace{\frac{d}{dm} f(X(s+m), s) \Big|_{m=0}}_{\text{'indirect effect'}} + \underbrace{f_2(X(s), s)}_{\text{'direct effect'}}$$

or (integrated & rearranged)

$$\int_r^t \frac{d}{dm} f(X(s+m), s) \Big|_{m=0} ds = V_X(t) - V_X(r) - \int_r^t f_2(X(s), s) ds.$$

'Outer' version is valid without classical assump'ns:

Identity lemma. If V_X is AC, then

$$\frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0} = V_X(t) - V_X(r) - \int_r^t f_2(X(s), s) ds.$$

(Where both sides are well-defined.)

(sketch proof: slide 47)

Proof of the main theorem

Identity lemma. If V_X is AC, then

$$\underbrace{\frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0}}_{\substack{= 0 \text{ for all } r, t \\ \text{iff oFOC holds.}}} = \underbrace{V_X(t) - V_X(r) - \int_r^t f_2(X(s), s) ds}_{\substack{= 0 \text{ for all } r, t \\ \text{iff } \boxtimes \text{ formula holds.}}}$$

Proof of main theorem. X satisfies the oFOC & V_X is AC
 \implies identity lemma applies. So oFOC \implies \boxtimes formula.

X satisfies the \boxtimes formula \implies V_X is AC (by Lebesgue's FToC)
 \implies identity lemma applies. So \boxtimes formula \implies oFOC. ■

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Environment

Agent with preferences $f(y, p, t)$ over physical outcome $y \in \mathcal{Y}$ and payment $p \in \mathbf{R}$.

- \mathcal{Y} partially ordered
- $f(y, \cdot, t)$ strictly decreasing and onto \mathbf{R}
- type $t \in [0, 1]$ is agent's private info.

Direct mechanism (Y, P) :

- *physical allocation* $Y : [0, 1] \rightarrow \mathcal{Y}$
- *payment rule* $P : [0, 1] \rightarrow \mathbf{R}$.

IC mechanism: $f(Y(t), P(t), t) \geq f(Y(r), P(r), t) \quad \forall r, t$.

Y is *implementable* iff \exists payment rule $P : [0, 1] \rightarrow \mathbf{R}$
s.t. (Y, P) is incentive-compatible.

Goal

‘Single-crossing’ preferences f :

higher types are more willing to pay to increase $y \in \mathcal{Y}$.

Theorem schema.

If \mathcal{Y} and f are ‘regular’ & f is ‘single-crossing’,
then any increasing allocation is implementable.

Versions:

- $\mathcal{Y} \subseteq \mathbf{R}$
 - classical assump’ns { Mirrlees (1976), Spence (1974),
Guesnerie and Laffont (1984)
 - no classical assump’ns Nöldeke and Samuelson (2018)
- general \mathcal{Y}
 - quasi-linear f { Matthews and Moore (1987),
García (2005)
 - general f this paper.

Regularity and single-crossing

Regular \mathcal{Y} : ‘rich’ & ‘not too large’. (def’n: slide 50)

Examples:

- \mathbf{R}^n ordered by ‘coordinate-wise smaller’
- finite-expectation RVs ordered by ‘a.s. smaller’
- distributions of posteriors updated from a given prior ordered by Blackwell.

Regular f : (def’n: slide 51)

- type derivative f_3 exists, bounded, continuous in p .
- f jointly continuous (when \mathcal{Y} has order topology).

Single-crossing f : (def’n: slide 52)

if type t willing to pay to increase $y \in \mathcal{Y}$, then so is type $t' > t$.

Implementability theorem.

If \mathcal{Y} and f are regular & f is single-crossing,
then any increasing physical allocation is implementable.

Argument I

Fix an increasing physical allocation $Y : [0, 1] \rightarrow \mathcal{Y}$.

Choose a payment rule P so that \boxtimes holds.

- \boxtimes formula is integral eq'n in unknown $P : [0, 1] \rightarrow \mathbf{R}$:

$$\phi(P(t), t) = \phi(P(0), 0) + \int_0^t \psi(P(s), s) ds$$

where $\phi(p, t) := f(Y(t), p, t)$ & $\psi(p, t) := f_3(Y(t), p, t)$

- f regular \implies solution exists.

Argument II

Embed into general setting:
$$\begin{cases} x := (y, p) \in \mathcal{Y} \times \mathbf{R} =: \mathcal{X} \\ f(x, t) = f(y, p, t) \\ \text{mechanisms } X = (Y, P). \end{cases}$$

f regular \implies type derivative exists & is bounded.

By *converse envelope theorem*, $\boxtimes \implies$ oFOC
 \iff mechanism (Y, P) is locally IC.

Finally, local IC \implies global IC by single-crossing.

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Selling information

Monopolist selling information.

Physical allocations \mathcal{Y} :
distributions of posterior beliefs, ordered by Blackwell.

Selling information: model

Finite set Ω of states. Prior $\mu_0 \in \text{int } \Delta(\Omega)$.

Distribution of posteriors: a Borel probability y on $\Delta(\Omega)$.

Fact: y is feasible iff $\int_{\Delta(\Omega)} \mu y(d\mu) = \mu_0$.

Physical outcomes \mathcal{Y} : dist'ns of posteriors with mean μ_0 .

Equip \mathcal{Y} with Blackwell order: $y \lesssim y'$ iff

$$\int_{\Delta(\Omega)} v dy \leq \int_{\Delta(\Omega)} v dy' \quad \forall \text{ cont's convex } v : \Delta(\Omega) \rightarrow \mathbf{R}.$$

\mathcal{Y} is regular.

Selling information: preferences

Agent's value of information: $V(\mu, t)$.

- V convex in μ
- V_2 exists, bounded, cont's in μ
- micro-foundation: $V(\mu, t) = \sup_{a \in A} \sum_{\omega \in \Omega} U(a, \omega, t) \mu(\omega)$.

Agent's preferences: $f(y, p, t) := g\left(\int_{\Delta(\Omega)} V(\mu, t) y(d\mu), p\right)$.

Assume single-crossing: higher types more willing
to pay for extra information.

Selling information: implementability

Information allocation: $Y : [0, 1] \rightarrow \mathcal{Y}$.

Implementability theorem \implies any Blackwell-increasing information allocation is implementable.

Thanks!



Absolute equi-continuity: properties

A family $\{\phi_x\}_{x \in \mathcal{X}}$ of functions $[0, 1] \rightarrow \mathbf{R}$ is *AEC* iff the family

$$\left\{ t \mapsto \sup_{x \in \mathcal{X}} \left| \frac{\phi_x(t+m) - \phi_x(t)}{m} \right| \right\}_{m>0} \text{ is uniformly integrable.}$$

Name inspired by the following (Fitzpatrick & Hunt, 2015):

AC–UI lemma. A continuous $\phi : [0, 1] \rightarrow \mathbf{R}$ is AC iff

$$\left\{ t \mapsto \frac{\phi(t+m) - \phi(t)}{m} \right\}_{m>0} \text{ is uniformly integrable.}$$

As name ‘AEC’ suggests, an AEC family

- is (uniformly) equi-continuous
- has AC functions as its members.

↔ back to slide 16

Sketch proof of the necessity lemma

Necessity lemma. Any optimal decision rule X satisfies the oFOC & has $V_X(t) := f(X(t), t)$ AC.

Sketch proof. X optimal & $\{f(x, \cdot)\}_{x \in \mathcal{X}}$ AEC $\implies V_X$ AC.

Since X optimal, have for any s and $m > 0 > m'$ that

$$\frac{f(X(s+m), s) - f(X(s), s)}{m} \leq 0 \leq \frac{f(X(s+m'), s) - f(X(s), s)}{m'}.$$

Integrating over (r, t) and letting $m, m' \rightarrow 0$,

both sides (in fact) converge to same limit:

$$\frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0} \leq 0 \leq \frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0}. \quad \blacksquare$$

\hookrightarrow back to slide 26

Sketch proof of the identity lemma I

$$\begin{aligned} \text{For } m > 0, \text{ write} \quad & \left. \frac{V_X(t+m) - V_X(t)}{m} \right\} =: \phi_m(t) \\ = \quad & \left. \frac{f(X(t+m), t+m) - f(X(t+m), t)}{m} \right\} =: \psi_m(t) \\ + \quad & \left. \frac{f(X(t+m), t) - f(X(t), t)}{m} \right\} =: \chi_m(t). \end{aligned}$$

$$\lim_{m \downarrow 0} \int_r^t \chi_m = \frac{d}{dm} \int_r^t f(X(s+m), s) ds \Big|_{m=0} \quad \text{if limit exists.}$$

Must show: limit exists & equals

$$V_X(t) - V_X(r) - \int_r^t f_2(X(s), s) ds.$$

Sketch proof of the identity lemma II

$$\begin{aligned} & \left. \frac{V_X(t+m) - V_X(t)}{m} \right\} =: \phi_m(t) \\ = & \left. \frac{f(X(t+m), t+m) - f(X(t+m), t)}{m} \right\} =: \psi_m(t) \\ + & \left. \frac{f(X(t+m), t) - f(X(t), t)}{m} \right\} =: \chi_m(t). \end{aligned}$$

$\{\psi_m\}_{m>0}$ need not converge a.e. (Even with strong assump'ns.)

But consider $\psi_m^*(t) := \frac{f(X(t), t) - f(X(t), t-m)}{m}$.

$\{\psi_m^*\}_{m>0}$ is UI & converges pointwise to $t \mapsto f_2(X(t), t)$. And

$$\begin{aligned} \int_r^t \psi_m &= \int_{r+m}^{t+m} \psi_m^* = \int_r^t \psi_m^* + \left(\int_t^{t+m} \psi_m^* - \int_r^{r+m} \psi_m^* \right) \\ &= \int_r^t \psi_m^* + o(1) \end{aligned} \quad \text{by UI.}$$

Sketch proof of the identity lemma III

$$\begin{aligned} & \left. \frac{V_X(t+m) - V_X(t)}{m} \right\} =: \phi_m(t) \\ = & \left. \frac{f(X(t+m), t+m) - f(X(t+m), t)}{m} \right\} =: \psi_m(t) \\ + & \left. \frac{f(X(t+m), t) - f(X(t), t)}{m} \right\} =: \chi_m(t). \end{aligned}$$

$$\int_r^t \psi_m = \int_r^t \psi_m^* + o(1), \quad \{\psi_m^*\}_{m>0} \text{ UI \& converges pointwise to } t \mapsto f_2(X(t), t).$$

V_X AC $\implies \{\phi_m\}_{m>0}$ UI & converges a.e. to V'_X . So

$$\begin{aligned} \lim_{m \downarrow 0} \int_r^t \chi_m &= \lim_{m \downarrow 0} \int_r^t [\phi_m - \psi_m] = \lim_{m \downarrow 0} \int_r^t [\phi_m - \psi_m^*] \\ \text{(Vitali)} \quad &= \int_r^t \lim_{m \downarrow 0} [\phi_m - \psi_m^*] = \int_r^t [V'_X(s) - f_2(X(s), s)] ds \\ \text{(FToC)} \quad &= V_X(t) - V_X(r) - \int_r^t f_2(X(s), s) ds. \end{aligned}$$

Outcome regularity

A set \mathcal{A} partially ordered by \lesssim is

- (1) *order-dense-in-itself* iff for any $a < a'$ in \mathcal{A} , there is a $b \in \mathcal{A}$ such that $a < b < a'$,
 - (2) *chain-separable* iff for each chain $C \subseteq \mathcal{A}$, there is a countable set $B \subseteq \mathcal{A}$ that is order-dense in C ,[‡]
 - (3) *countably chain-complete* iff every countable chain in \mathcal{A} with a lower (upper) bound in \mathcal{A} has an infimum (a supremum) in \mathcal{A} .
- (1) & (2): \mathcal{A} ‘rich’. (3): \mathcal{A} ‘not too large’.

Definition. \mathcal{Y} is *regular* iff it satisfies properties (1)–(3).

↔ back to slide 35

[‡] $B \subseteq \mathcal{A}$ is *order-dense* iff for any $a < a'$ in \mathcal{A} , $\exists b \in B$ s.t. $a \lesssim b \lesssim a'$.

Preference regularity

Order topology on a set \mathcal{A} partially ordered by \lesssim :
the topology generated by the open order rays

$$\{b \in \mathcal{A} : b < a\} \quad \text{and} \quad \{b \in \mathcal{A} : a < b\}.$$

Definition. f is *regular* iff

- (a) type derivative f_3 exists & is bounded & continuous in p
- (b) for any chain $\mathcal{C} \subseteq \mathcal{Y}$, f jointly continuous on $\mathcal{C} \times \mathbf{R} \times [0, 1]$
when \mathcal{C} has relative top'gy inherited from order top'gy on \mathcal{Y} .

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Single-crossing

Definition. For $\phi : [0, 1] \rightarrow \mathbf{R}$, upper & lower derivatives

$$D^* \phi(t) := \limsup_{m \rightarrow 0} \frac{\phi(t+m) - \phi(t)}{m}$$

$$D_* \phi(t) := \liminf_{m \rightarrow 0} \frac{\phi(t+m) - \phi(t)}{m}.$$

Partial upper/lower derivatives: $(D^*)_i$ & $(D_*)_i$.

Definition. f is *single-crossing* iff

for any increasing $Y : [0, 1] \rightarrow \mathcal{Y}$ & any $P : [0, 1] \rightarrow \mathbf{R}$,
mis-reporting payoff $U(r, t) := f(Y(r), P(r), t)$ satisfies

$$(D^*)_1 U(t, t) \geq 0 \quad \text{implies} \quad (D_*)_1 U(t, t') > 0 \quad \text{for } t' > t$$

$$\text{and } (D_*)_1 U(t, t) \leq 0 \quad \text{implies} \quad (D^*)_1 U(t, t') < 0 \quad \text{for } t' < t.$$

↔ back to slide 35

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